

# Example Induction Proofs on Trees

CS 173 Lecture B Fall 2016

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# Induction proof on trees

**Claim:** In a binary tree of height  $h$ , the number of nodes  $n \leq 2^{h+1} - 1$ .

**Proof:** Use Induction on height of tree.

Base case:  $h = 0$

$$2^{h+1} - 1 = 2^1 - 1 = 1$$

$$n = 1$$

$$\therefore n \leq 2^{h+1} - 1$$

(cont'd.)

$\textcircled{1}$   $T$   $h = 0$   
Tree  
can only  
have  
one  
node

Claim: In a binary tree of height  $h$ , the number of nodes  $n \leq 2^{h+1} - 1$ .

**Proof (continued):**

**Induction Hypothesis:**  $\forall h \leq k, n \leq 2^{h+1} - 1$ .

Want to prove: For  $h = k + 1, n \leq 2^{h+2} - 1$ .

**Induction step:** Recall that in each node in a binary tree has at most two children.

Consider a tree with height  $k + 1$  with  $n$  nodes. Let's call its root node  $r_0$ . There are two possibilities:  $r_0$  has one child, or  $r_0$  has two children.

**Case 1:**  $r_0$  has (exactly) one child. Call the child  $r_1$ .

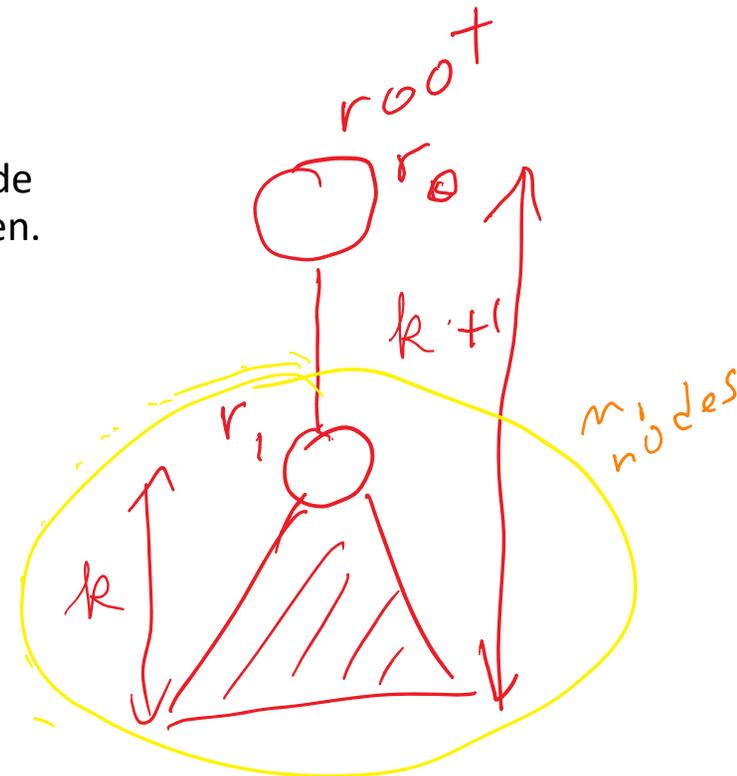
Suppose the subtree with root  $r_1$  has  $n_1$  nodes. Then  $n = n_1 + 1$ .

Now the height of the subtree rooted at  $r_1$  must be  $k$  (otherwise the height of the entire tree rooted at  $r_0$  could not be  $k + 1$ ).

By induction hypothesis:  $n_1 \leq 2^{k+1} - 1$ . Therefore  $n \leq 2^{k+1}$ .

Now  $2^{k+1} \leq 2^{k+2} - 1$  for all  $k \geq 0$ , thus  $n \leq 2^{h+2} - 1$ .

We proceed to Case 2.



Claim: In a binary tree of height  $h$ , the number of nodes  $n \leq 2^{h+1} - 1$ .

Proof (continued):

Induction Hypothesis:  $\forall h \leq k, n \leq 2^{h+1} - 1$ .

Want to prove: For  $h = k + 1, n \leq 2^{k+2} - 1$ .

**Proof (continued). Case 2:** Suppose tree rooted at  $r_0$  has height  $k+1$  and (exactly) two children. Call the children  $r_1$  and  $r_2$ .

Suppose the subtrees rooted at  $r_1$  and  $r_2$  have  $n_1$  and  $n_2$  nodes, and heights  $k_1$  and  $k_2$ , respectively.

Now  $n = n_1 + n_2 + 1, k = \max(k_1, k_2) + 1, k_1 \leq k$ , and  $k_2 \leq k$ .

Without loss of generality suppose  $k_1 \leq k_2$ . This implies  $k = k_2$ .

By induction hypothesis:

$$n_1 \leq 2^{k_1+1} - 1 \text{ and } n_2 \leq 2^{k_2+1} - 1$$

$$\text{Thus } n \leq (2^{k_1+1} - 1) + (2^{k_2+1} - 1) + 1,$$

$$\text{i.e. } n \leq 2^{k_2+1} + 2^{k_2+1} - 1. \text{ So } n \leq 2^{k_2+2} - 1. \text{ But } k = k_2.$$

$$\text{This implies } n \leq 2^{k+2} - 1.$$

Which is what we wanted to prove. QED.

